

Universal Turing Machines

Lecture 30
Section 10.4

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1 Universal Turing Machines

2 Enumerators

3 Assignment

Outline

- 1 Universal Turing Machines
- 2 Enumerators
- 3 Assignment

Universal Turing Machines

Definition (Universal Turing Machines)

A **universal Turing machine** is a Turing machine that can simulate any Turing machine.

- The input tape contains two items:
 - An encoding of the Turing machine M to be simulated.
 - The input w to be read by M .
- The universal Turing machine U will read M and w and
 - Accept, reject, or loop, according to what M would do when reading w ,
 - Write to the tape the output that M would write when reading w .

Universal Turing Machines

- The encoding is straightforward:
 - The states $\{q_1, q_2, \dots, q_n\}$ are encoded as $1, 11, \dots, 111 \dots 1$ and separated by 0's.
 - q_1 is the start state.
 - q_2 is the sole final state.
 - The tape symbols $\{a_1, a_2, \dots, a_m\}$ are encoded as $1, 11, \dots, 111 \dots 1$ and separated by 0's.
 - “Left” is encoded as 1.
 - “Right” is encoded as 11.
- For example, the transition $\delta(q_2, a_3) = (q_1, a_4, R)$ would be encoded as

1101110101111011

Universal Turing Machines

- The universal Turing machine U has three tapes.
 - Tape 1 contains the encoding of a machine M .
 - Tape 2 contains the encoded input to M .
 - Tape 3 contains the current state.

Universal Turing Machines

- U reads the current state q_i from Tape 3 and the current symbol a_j from Tape 2.
- It searches Tape 1 for a transition $(q_i, a_j) = (q_k, a_m, d)$.
- When it finds it, it
 - Replaces q_i with q_k on Tape 3.
 - Replaces a_j with a_m on Tape 2.
 - Moves left or right on Tape 2, according to d .
- If and when it fails to find (q_i, a_j) on Tape 1, it quits.

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Enumeration

Definition (Enumerate)

To **enumerate** a language L is to list all of the strings in L , each string listed exactly once. The order does not matter.

Definition (Enumerator)

An **enumerator** for a language L is a Turing machine that enumerates L on its tape.

- If L is infinite, then clearly the enumerator does not halt.

Enumerators

Example (Enumerator)

- Let $\Sigma = \{0, 1\}$.
- It is easy to build an enumerator for Σ^* .
 - Begin with \$ on the tape.
 - Write λ “on the tape.”
 - Move right, write a separator \$, and write 0.
 - Copy the last number written and increment the copy, unless the copy is all 1’s, in which case replace it with all 0’s and one additional 0.
 - Write the separator \$.
 - Repeat the previous three steps.
- The tape contents will be

$\$ \$ 0 \$ 1 \$ 0 0 \$ 0 1 \$ 1 0 \$ 1 1 \$ 0 0 0 \dots$

Enumerators

- Let $\Sigma = \{0, 1\}$.
- Describe enumerators for the following languages.
 - All strings with an equal number of 0's and 1's.
 - All legitimate patterns of parentheses, where 0 represents '(' and 1 represents ')'.
 - All prime numbers.
 - All numbers that divide one of the decimal integers 9, 99, 999, 9999, 99999, ...
 - All pairs $\{a, b\}$, where $a, b \geq 0$.
 - All triples $\{a, b, c\}$, where $a, b, c \geq 0$.

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Assignment

Homework

- Section 10.3 Exercises 1, 2, 4, 5, 6, 8, 9, 10.